

Rhythmic Syncope and Strict Locality in Subregular Phonology

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Introduction





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Rhythmic Syncope

Definition

Assume an alphabet $\Sigma = C \cup V$, with $C \cap V = \emptyset$. The *rhythmic* syncope function is the function defined by

 $\rho(\boldsymbol{c}_0\boldsymbol{v}_1\boldsymbol{c}_1\boldsymbol{v}_2\boldsymbol{c}_2\ldots\boldsymbol{v}_n\boldsymbol{c}_n) = \boldsymbol{c}_0\boldsymbol{v}_1\boldsymbol{c}_1\boldsymbol{c}_2\boldsymbol{v}_3\boldsymbol{c}_3\boldsymbol{c}_4\ldots\boldsymbol{c}_n$

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where for each *i*, $c_i \in C^*$ and $v_i \in V$.

 $\rho(\textit{CVCVCVCVCVC}) = \textit{CVCCVCCVC}$



Outline

• Show that ρ is not *strictly local*.



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- Define a class of functions that includes ρ .
- Discuss the theoretical consequences.



Strictly Local Functions

Subregular Phonology



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- ▶ Vowel Harmony: *ea, *eı, *eu, *ae, *ai, *aü...



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- ► x ∈ L if and only if no element of S is a substring of x, ignoring symbols not in T.



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► Parsimony: This common prefix is the longest one possible.



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For a function $f: \Sigma^* \to \Sigma^*$,

$$f^{\leftarrow}(\textbf{\textit{x}}) = \mathsf{lcp}(\{f(\textbf{\textit{xz}}) | \textbf{\textit{z}} \in \Sigma^*\})$$

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Definition

For a string $x \in \Sigma^*$, tier $T \subseteq \Sigma$, and number k, suff^k_T(x) is the last k symbols of x on tier T.



Definition (Chandlee et al., In prep)

A function $f: \Sigma^* \to \Sigma^*$ is *k*-strictly local on tier *T* if for every $u, v \in \Sigma^*$, if $\operatorname{suff}_{\tau}^{k-1}(u) = \operatorname{suff}_{\tau}^{k-1}(v)$

and

$$\operatorname{suff}_{T}^{k-1}(f^{\leftarrow}(u)) = \operatorname{suff}_{T}^{k-1}(f^{\leftarrow}(v)),$$

then for all $w \in \Sigma^*$ we have

$$f^{\rightarrow}(w, uw) = f^{\rightarrow}(w, vw).$$



To show that ρ is *not* k-SL on tier *T*, we must find u, v, w such that

$$\begin{split} & \mathsf{suff}_{\mathsf{T}}^{k-1}(u) = \mathsf{suff}_{\mathsf{T}}^{k-1}(v) \\ & \mathsf{suff}_{\mathsf{T}}^{k-1}(\rho^\leftarrow(u)) = \mathsf{suff}_{\mathsf{T}}^{k-1}(\rho^\leftarrow(v)), \end{split}$$

 $\mathsf{but}\,\rho^{\rightarrow}(\mathbf{\textit{w}},\mathbf{\textit{uw}})\neq\rho^{\rightarrow}(\mathbf{\textit{w}},\mathbf{\textit{vw}}).$



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$$\rho^{\leftarrow}(\mathbf{u}) = \mathbf{a}^{2k}, \rho^{\leftarrow}(\mathbf{v}) = \mathbf{a}^{2k+1}$$

•
$$\rho(\mathbf{u}\mathbf{w}) = \mathbf{a}^{2\mathbf{k}+1}$$
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but $\rho^{\rightarrow}(w, uw) \neq \rho^{\rightarrow}(w, vw)$. • Let *a* be a vowel. $u = a^{4k}$, $v = a^{4k+1}$, w = a• $\rho^{\leftarrow}(u) = a^{2k}$, $\rho^{\leftarrow}(v) = a^{2k+1}$ • $\rho(uw) = a^{2k+1}$, $\rho(vw) = a^{2k+1}$ • $\rho^{\rightarrow}(w, uw) = a$, $\rho^{\rightarrow}(w, vw) = \emptyset$



Time Alignment







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Deletion destroys evidence.





Existing OT analyses address this problem.

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Time-Aligned TSL Functions

Definition

Let $f: \Sigma^* \to \Sigma^*$ and $x = x_1 x_2 \dots x_n \in \Sigma^*$. The *i*th most recent action of f on x is the pair $\langle x_{n-i+1}, f_i^{\leftarrow}(x) \rangle$, where $f_i^{\leftarrow}(x)$ is the string such that

$$f^{\leftarrow}(\mathbf{x}_1\mathbf{x}_2\ldots\mathbf{x}_{n-i+1}) = f^{\leftarrow}(\mathbf{x}_1\mathbf{x}_2\ldots\mathbf{x}_{n-i})f_i^{\leftarrow}(\mathbf{x}).$$

For $T \subseteq \Sigma^*$, the *ith most recent action of f on x on tier T* is the action denoted

$$\langle \mathbf{x}_{i,T}, \mathbf{f}_{i,T}^{\leftarrow}(\mathbf{x}) \rangle := \langle \mathbf{x}_{n-j+1}, \mathbf{f}_{n-j+1}^{\leftarrow}(\mathbf{x}) \rangle,$$

where *j* is the *i*th largest index such that $x_{n-j+1} \in T$ and $f_{n-j+1}^{\leftarrow}(x) \in T^*$.



Time-Aligned TSL Functions

Definition

Let $f : \Sigma^* \to \Sigma^*$ and $T \subseteq \Sigma$. For $k \in \mathbb{N}$, f is time-aligned k-strictly local on tier T if for all $u, v \in \Sigma$, if

$$\langle u_{i,T}, f_{i,T}^{\leftarrow}(u) \rangle = \langle v_{i,T}, f_{i,T}^{\leftarrow}(V) \rangle$$

for $1 \leq i \leq k$, then for all $w \in \Sigma^*$,

$$\boldsymbol{f}^{\rightarrow}(\boldsymbol{w},\boldsymbol{u}\boldsymbol{w}) = \boldsymbol{f}^{\rightarrow}(\boldsymbol{w},\boldsymbol{v}\boldsymbol{w}).$$





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$$\ \ \, \rho_{1,\mathbf{V}}^{\leftarrow}(\mathbf{c}_0\mathbf{v}_1\mathbf{c}_1\mathbf{v}_2\mathbf{c}_2\ldots\mathbf{v}_{2m}\mathbf{c}_{2m}) = \varnothing$$



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- Modern speakers understand the 1930s forms but do not use them.



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 - ► Today: $/gtigmmibna:-d/ \rightarrow [gtigmmibna:-d]$
- Modern speakers understand the 1930s forms but do not use them.
- Similar phenomena have been observed in Old Russian (Isacenko, 1970) and Old Irish (McManus, 1983).



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- Rhythmic syncope is not TSL.
- ► The time-aligned TSL functions incorporate rhythmic syncope.
- Rejection of rhythmic syncope by child learners would constitute evidence for the TSL hypothesis.



References

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